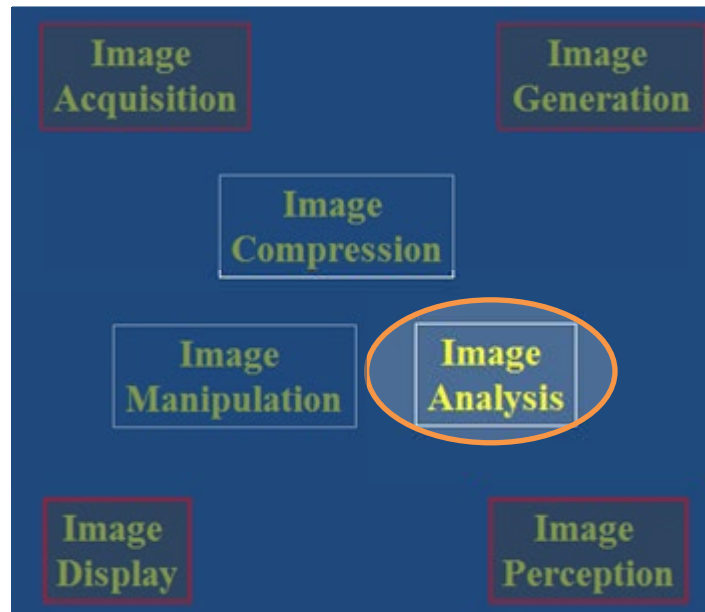


Digital Image Processing

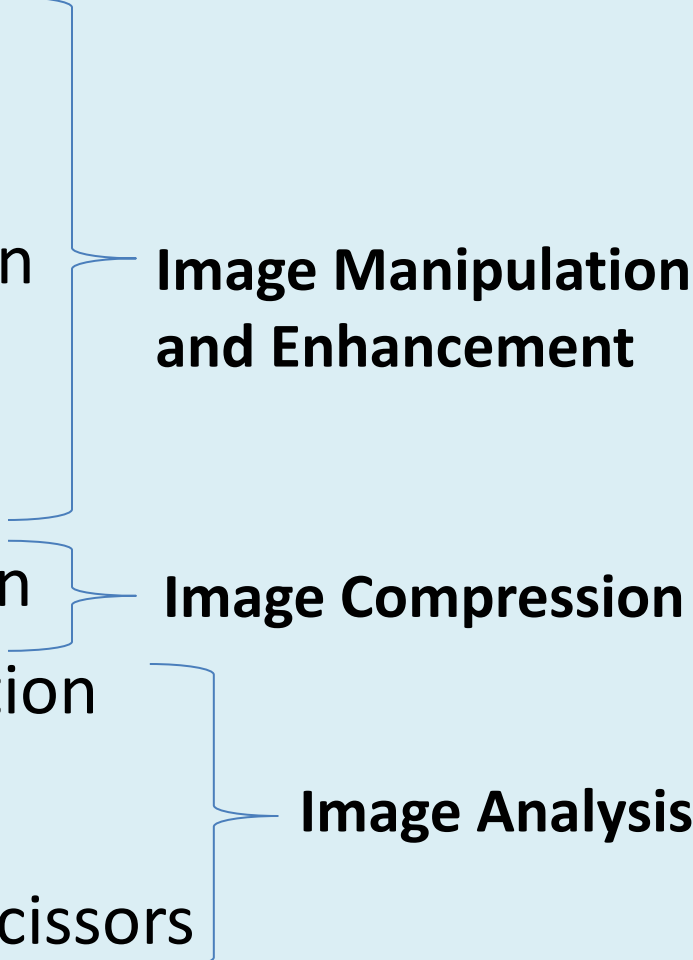
Edge Detection: Smart Scissors



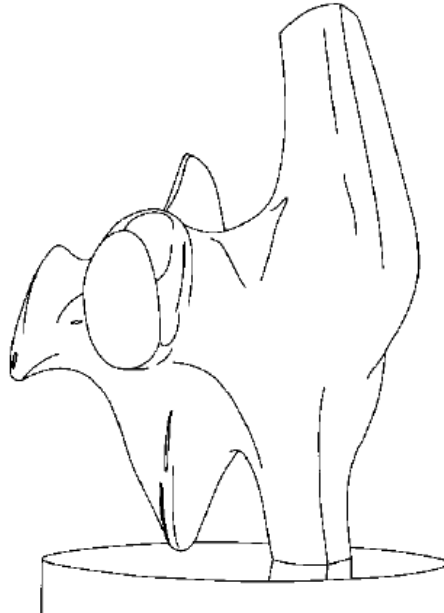
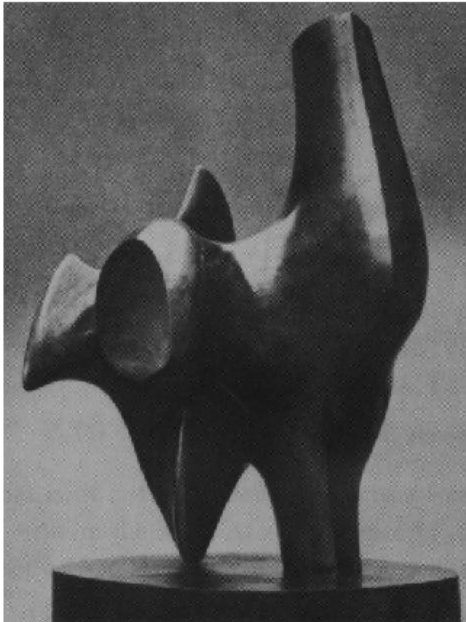
DRAFT



Lecture Objectives

- Previously
 - Filtering
 - Interpolation
 - Warping
 - Morphing
 - Compression
 - Edge Detection
 - Today
 - Intelligent Scissors
- Image Manipulation and Enhancement**
- Image Compression**
- Image Analysis**
- 

Recall: Edge Detection



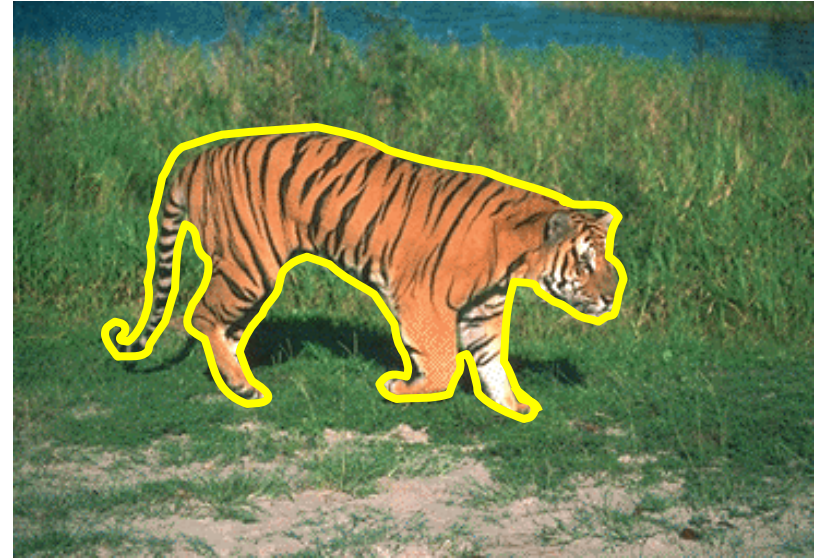
- **Goal:** Identify sudden changes (discontinuities) in an image
 - Much of the 'shape' information of the image can be realized in the edges
- **Ideal Result:** an artist's line drawing
 - *aside: who can use knowledge beyond a single image*

Recall: Canny Edge Detector



- Filter image with derivative of Gaussian
- Find magnitude and orientation of gradient
- Apply non-maximum suppression
- Linking and thresholding (hysteresis)
 - Define 2 thresholds: low and high
 - Use the high threshold to start edge curves and the low threshold to continue them

Extracting Objects



- How can the tiger be extracted from the image?
 - Difficult to do manually
 - Difficult to do automatically (but possible)
 - Easier to do semi-automatically

Paper to Read/Reference

- Intelligent Scissors,
Mortensen et. al, SIGGRAPH 1995

Intelligent Scissors

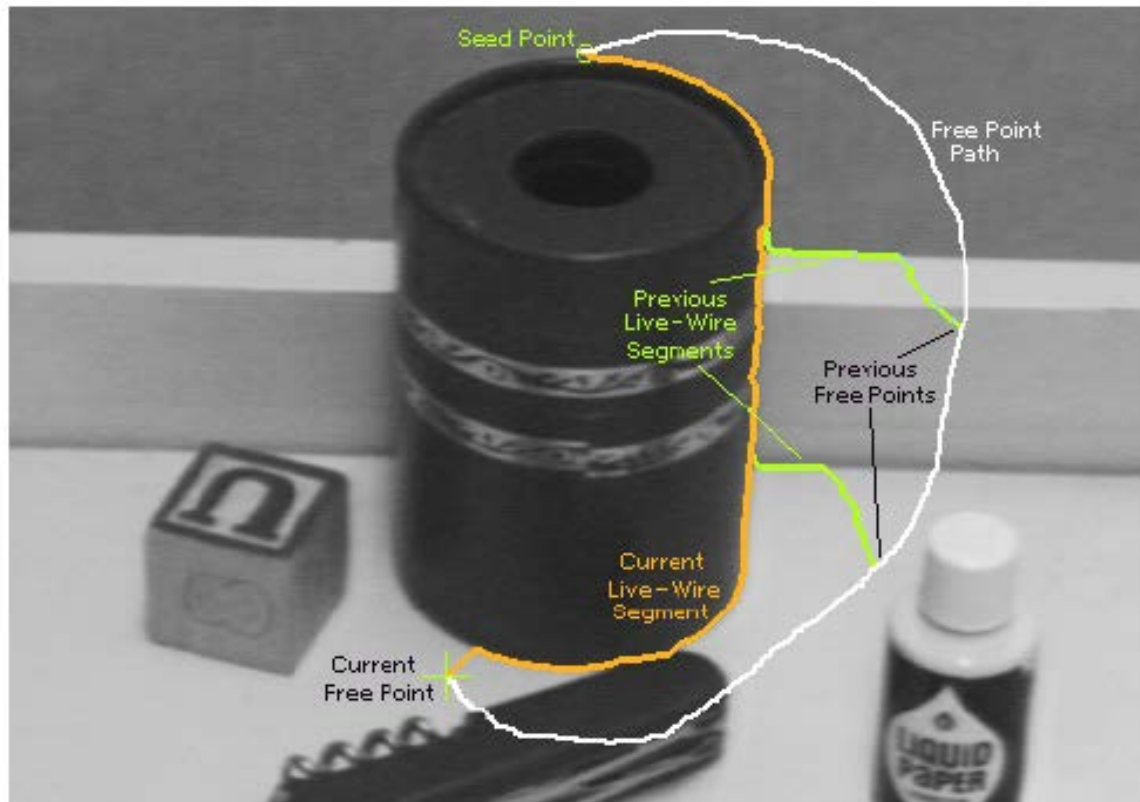
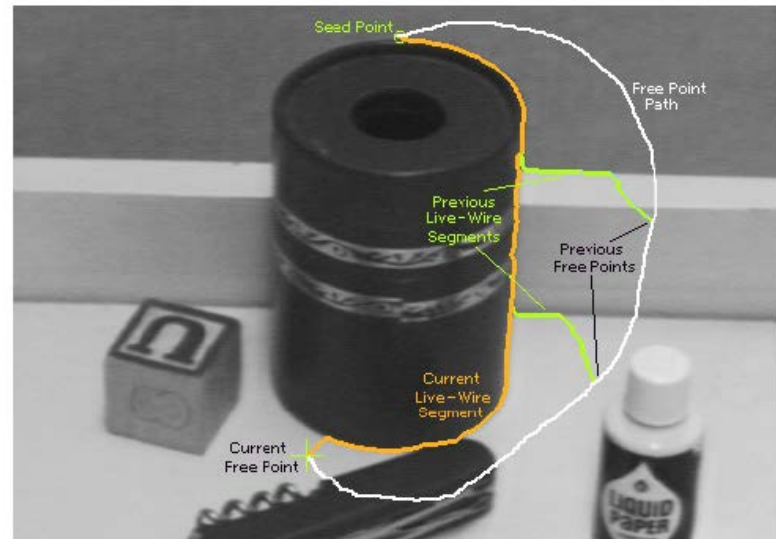


Figure 2: Image demonstrating how the live-wire segment adapts and snaps to an object boundary as the free point moves (via cursor movement). The path of the free point is shown in white. Live-wire segments from previous free point positions (t_0 , t_1 , and t_2) are shown in green.

Intelligent Scissors

- **Question**
 - How do we find a path from seed to mouse that follows an object boundary as closely as possible?
- **Answer**
 - Define a path that stays as close as possible to edges



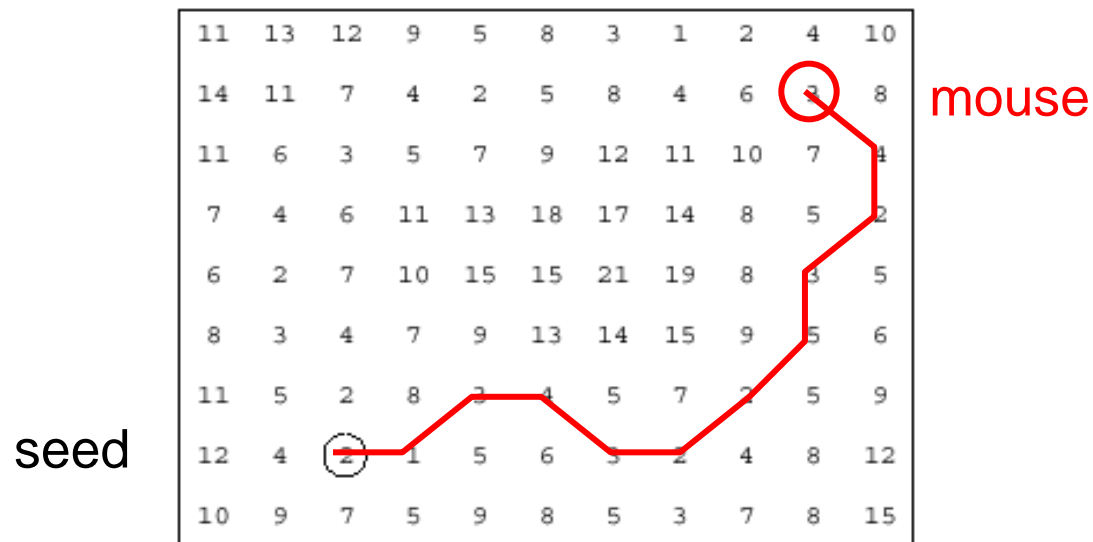
Intelligent Scissors

- Basic Idea

- Define edge score for each pixel

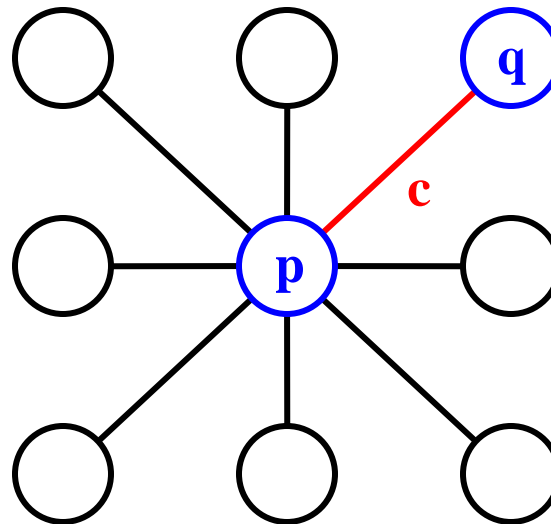
- assign edge pixels a low score

- Find lowest cost path from seed to mouse



Let's look at this more closely

- Treat the image as a graph



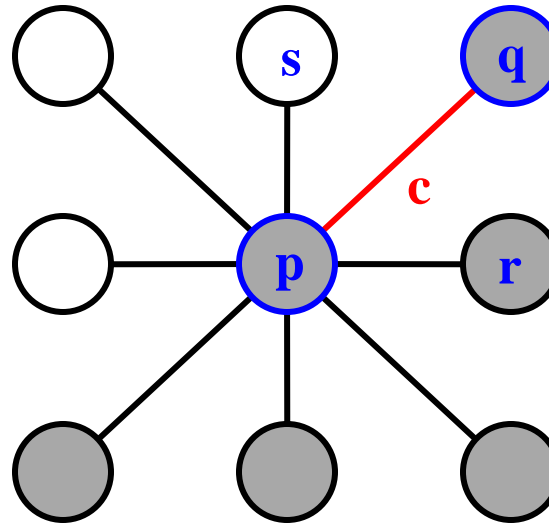
Graph

- node for every pixel **p**
- link between every adjacent pair of pixels, **p,q**
- cost **c** for each link

Note: each *link* has a cost

- this is a little different than the figure before where each pixel had a cost

Defining the costs

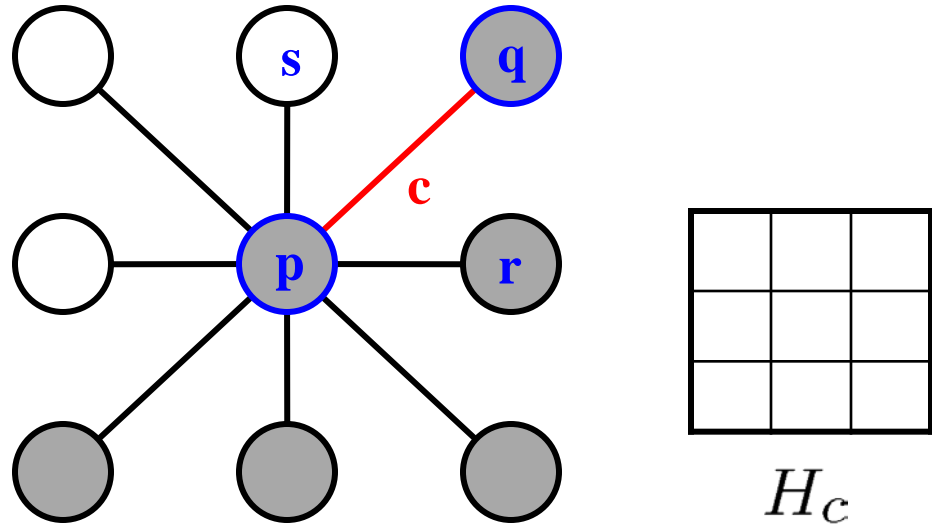


Want to hug image edges: how to define cost of a link?

- good (low-cost) links follow the intensity edge
 - want intensity to change rapidly \perp to the link

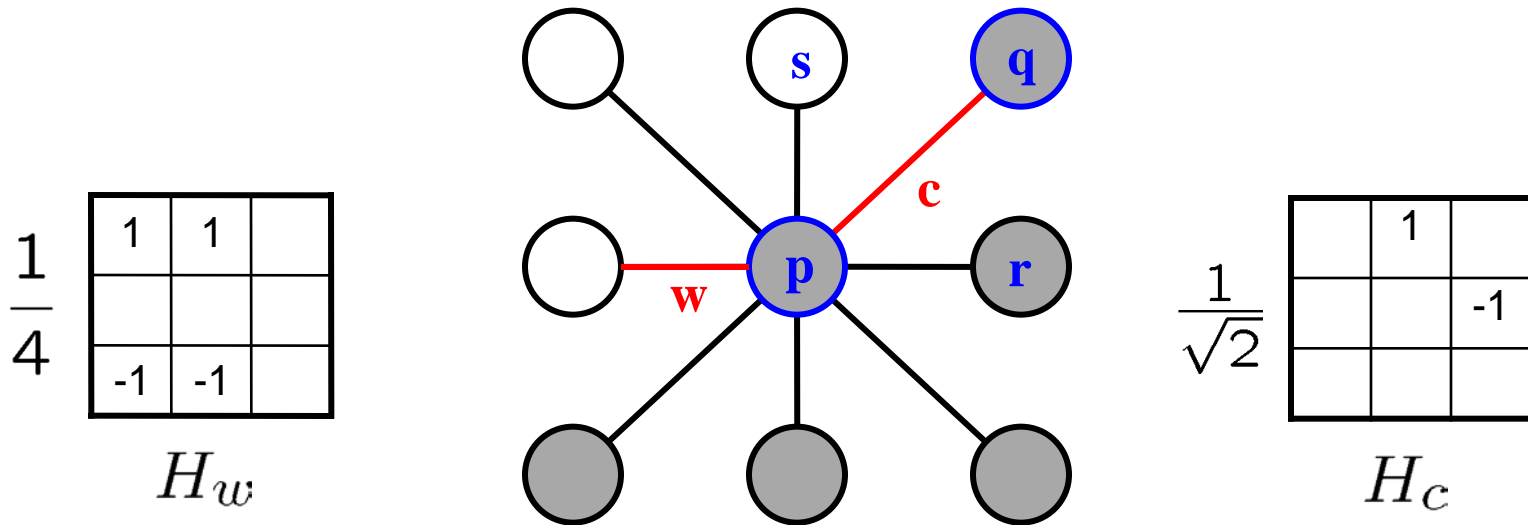
- $c \approx -\frac{1}{\sqrt{2}} |\text{intensity of } r - \text{intensity of } s|$

Defining the costs



- c can be computed using a cross-correlation filter
 - assume it is centered at p

Defining the costs



c can be computed using a cross-correlation filter

- assume it is centered at p

A couple more modifications

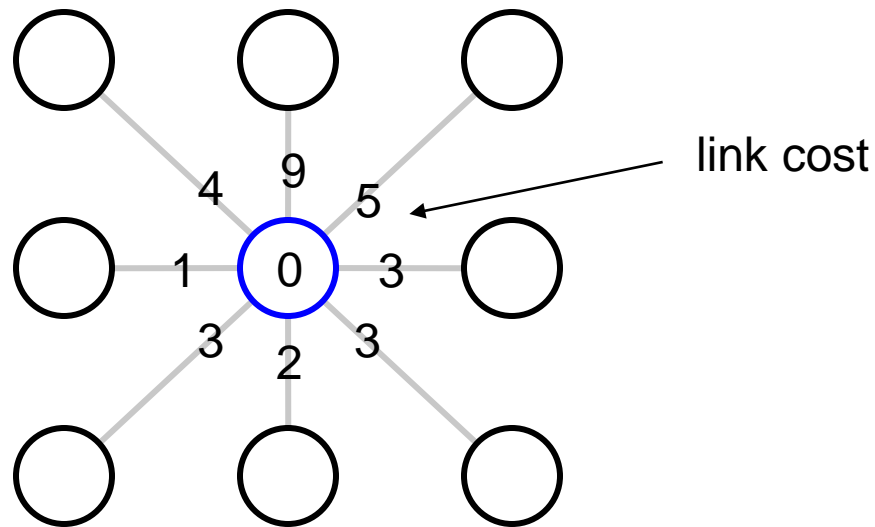
- Scale the filter response by length of link c . Why?

- Make c positive

– Set $c = (\max|\text{filter response}| \cdot \text{length})$

– where $\max = \text{maximum } |\text{filter response}| \cdot \text{length}$ over all pixels in the image source: Noah Snavely

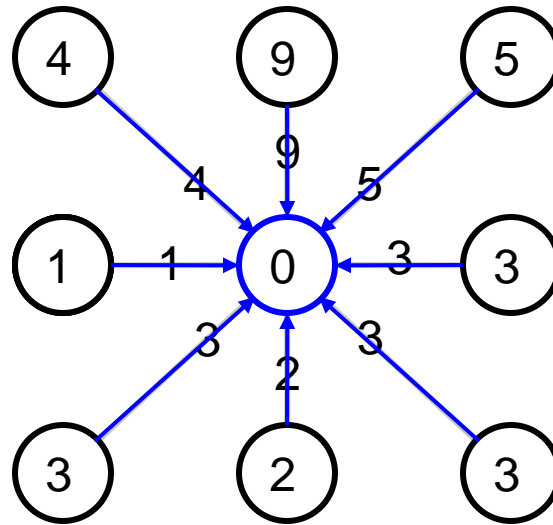
Dijkstra's shortest path algorithm



Algorithm

1. init node costs to ∞ , set p = seed point, $\text{cost}(p) = 0$
2. expand p as follows:
 - for each of p 's neighbors q that are not expanded
 - set $\text{cost}(q) = \min(\text{cost}(p) + c_{pq}, \text{cost}(q))$

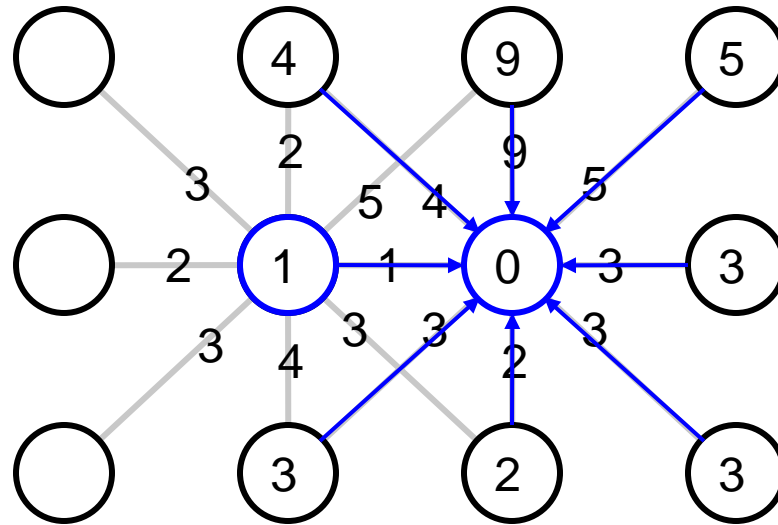
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 - if q 's cost changed, make q point back to p
 - put q on the ACTIVE list (if not already there)

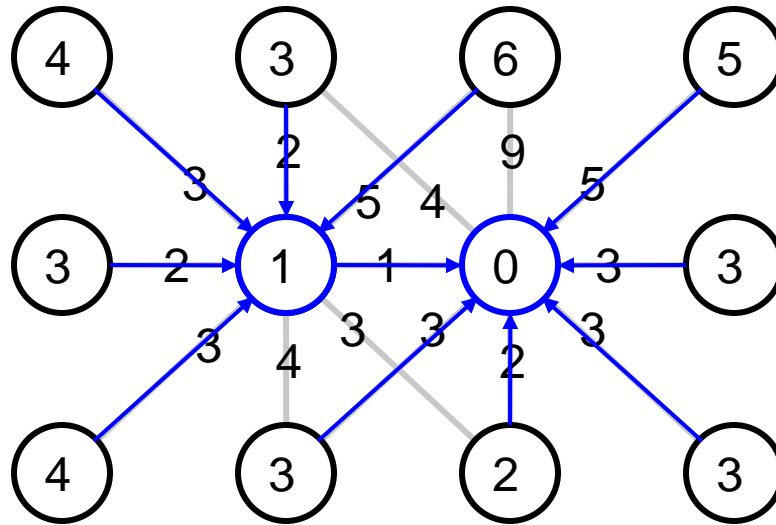
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3. set r = node with minimum cost on the ACTIVE list
4. repeat Step 2 for $p = r$

Dijkstra's shortest path algorithm



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Dijkstra's shortest path algorithm

- Properties

- It computes the minimum cost path from the seed to every node in the graph. This set of minimum paths is represented as a *tree*
- Running time, with N pixels:
 - $O(N^2)$ time if you use an active list
 - $O(N \log N)$ if you use an active priority queue (heap)
 - takes fraction of a second for a typical (640x480) image
- Once this tree is computed once, we can extract the optimal path from any point to the seed in $O(N)$ time.
 - it runs in real time as the mouse moves
- What happens when the user specifies a new seed?

Questions?

- Beyond D2L
 - Examples and information can be found online at:
 - *<http://docdingle.com/teaching/cs.html>*

- *Continue to more stuff as needed*

Extra Reference Stuff Follows

Credits

- Much of the content derived/based on slides for use with the book:
 - *Digital Image Processing*, Gonzalez and Woods
- Some layout and presentation style derived/based on presentations by
 - Donald House, Texas A&M University, 1999
 - Sventlana Lazebnik, UNC, 2010
 - Noah Snavely, Cornell University, 2012
 - Xin Li, WVU, 2014
 - George Wolberg, City College of New York, 2015
 - Yao Wang and Zhu Liu, NYU-Poly, 2015

